

PROGRAMMING FOR PROBLEM SOLVING -- PPS SUBJECT CODE- BTPS-101-18



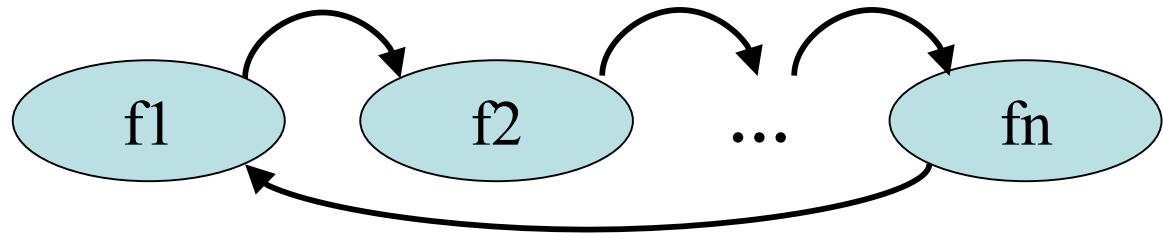
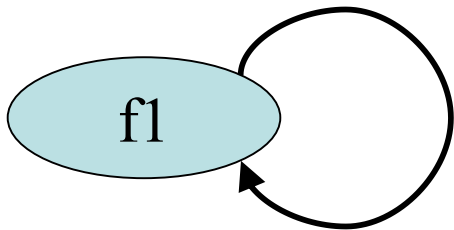
PPS

Unit-6 Recursion



Recursive Function

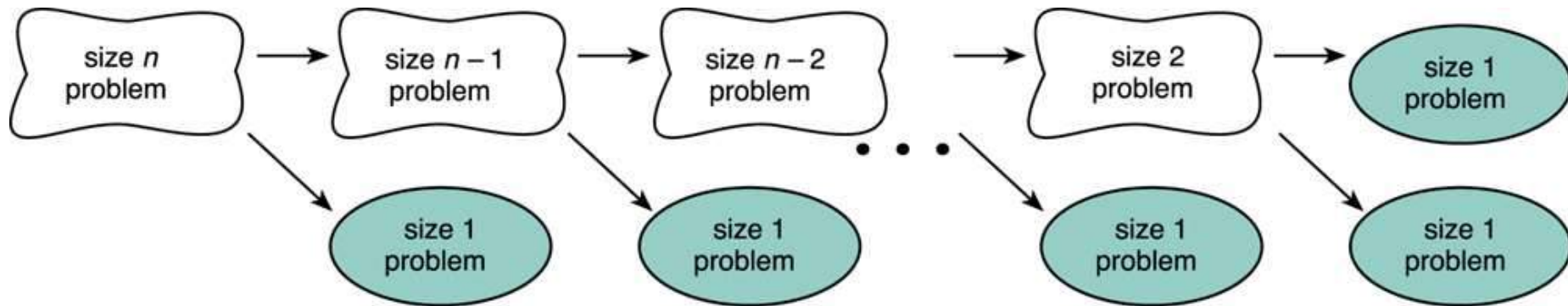
- The **recursive function** is
 - a kind of function that calls itself, or
 - a function that is part of a cycle in the sequence of function calls.



Problems Suitable for Recursive Functions

- One or more simple cases of the problem have a straightforward solution.
- The other cases can be redefined in terms of problems that are closer to the simple cases.
- The problem can be reduced entirely to simple cases by calling the recursive function.
 - *If this is a simple case
solve it
else
redefine the problem using recursion*

Splitting a Problem into Smaller Problems



- Assume that the problem of size 1 can be solved easily (i.e., the simple case).
- We can recursively split the problem into a problem of size 1 and another problem of size $n-1$.

An Example of Recursive Function

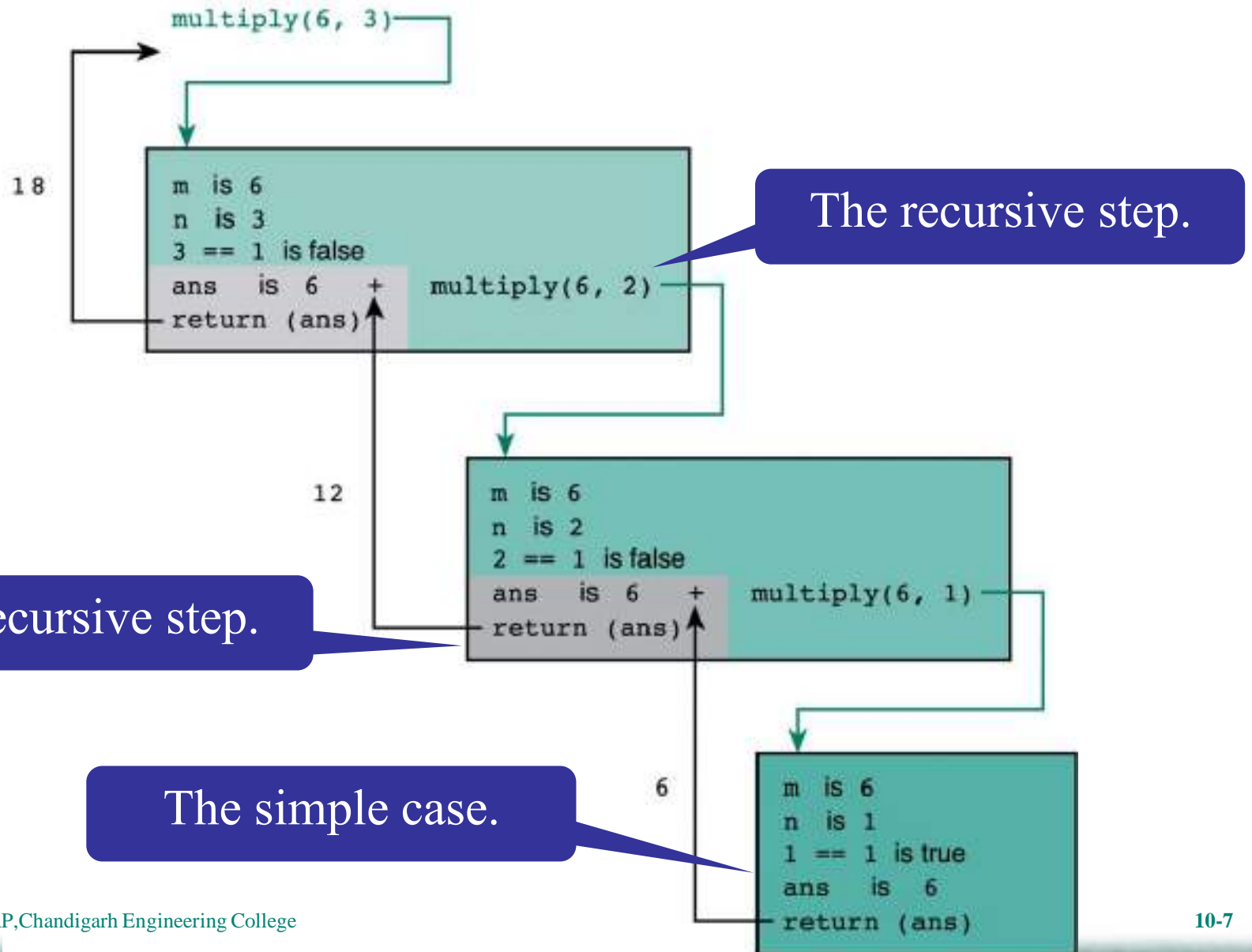
- We can implement the multiplication by addition.

```
1.  /*
2.  * Performs integer multiplication using + operator.
3.  * Pre:  m and n are defined and n > 0
4.  * Post: returns m * n
5.  */
6.  int
7.  multiply(int m, int n)
8.  {
9.      int ans;
10.
11.     if (n == 1)
12.         ans = m;      /* simple case */
13.     else
14.         ans = m + multiply(m, n - 1); /* recursive step */
15.
16.     return (ans);
17. }
```

The simple case is “ $m * 1 = m$.”

The recursive step uses the following equation: “ $m * n = m + m * (n - 1)$.”

Trace of Function multiply(6, 3)



Terminating Condition

- The recursive functions always contains one or more **terminating conditions**.
 - A condition when a recursive function is processing a simple case instead of processing recursion.
- Without the terminating condition, the recursive function may run forever.
 - e.g., in the previous multiply function, the if statement “`if (n == 1) ...`” is the terminating condition.

Recursive Function to Count a Character in a String

- We can count the number of occurrences of a given character in a string.
 - e.g., the number of 's' in "Mississippi" is 4.

```
4. int
5. count(char ch, const char *str)
6. {
7.
8.     int ans;
9.
10.    if (str[0] == '\0')                /* simple case */
11.        ans = 0;
12.    else                                /* redefine problem using recursion */
13.        if (ch == str[0])              /* first character must be counted */
14.            ans = 1 + count(ch, &str[1]);
15.        else                            /* first character is not counted */
16.            ans = count(ch, &str[1]);
17.
18.    return (ans);
19. }
```

The terminating condition.

A Recursive Function that Reverses Input Words (1/2)

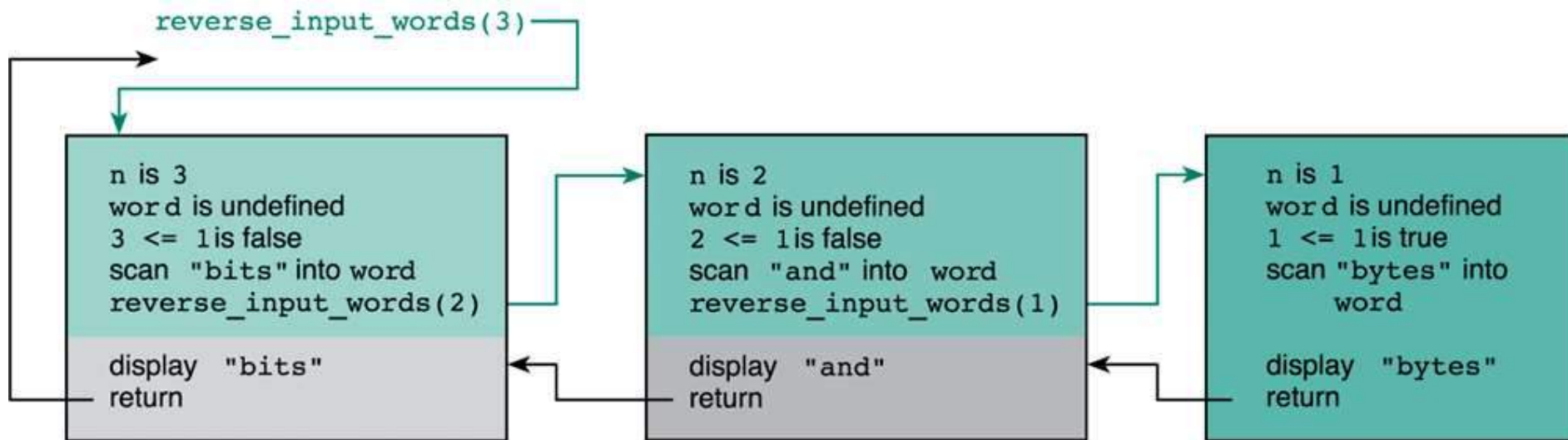
- The recursive concept can be used to reverse an input string.
 - It can also be done without recursion.

```
6. reverse_input_words(int n)
7. {
8.     char word[WORDSIZ]; /* local variable for storing one word */
9.
10.    if (n <= 1) { /* simple case: just one word to get and print */
11.
12.        scanf("%s", word);
13.        printf("%s\n", word);
14.
15.    } else { /* get this word;
16.            reverse order;
17.
18.            scanf("%s", word);
19.            reverse_input_words(n - 1);
20.            printf("%s\n", word);
21.
22.    }
```

The scanned word will not be printed until the recursion finishes.

The first scanned word is last printed.

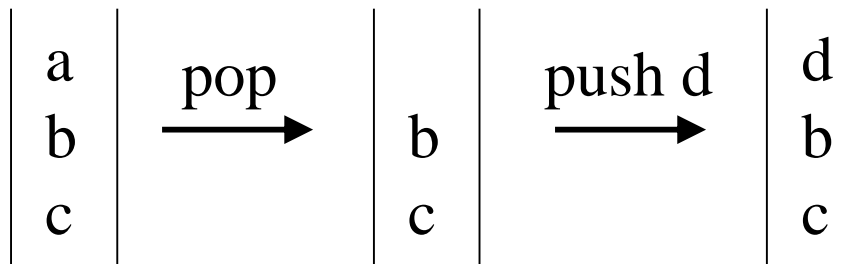
A Recursive Function that Reverses Input Words (2/2)



- Note that the recursive function is just an alternative solution to a problem.
 - You can always solve the problem without recursion.

How C Maintains the Recursive Steps

- C keeps track of the values of variables by the **stack** data structure.
 - Recall that stack is a data structure where the last item added is the first item processed.
 - There are two operations (push and pop) associated with stack.



How C Maintains the Recursive Steps

- Each time a function is called, the execution state of the caller function (e.g., parameters, local variables, and memory address) are pushed onto the stack.
- When the execution of the called function is finished, the execution can be restored by popping up the execution state from the stack.
- This is sufficient to maintain the execution of the recursive function.
 - The execution state of each recursive step are stored and kept in order in the stack.

How to Trace Recursive Functions

- The recursive function is not easy to trace and to debug.
 - If there are hundreds of recursive steps, it is not useful to set the breaking point or to trace step-by-step.
- A naïve but useful approach is inserting printing statements and then watching the output to trace the recursive steps.

```
7. int
8. multiply(int m, int n)
9. {
10.     int ans;
11.
12.     printf("Entering multiply with m = %d, n = %d\n", m, n);
13.
14.     if (n == 1)
15.         ans = m;      /* simple case */
16.     else
17.         ans = m + multiply(m, n - 1); /* recursive step */
```

Watch the input arguments passed into each recursive step.

(continued)

Recursive factorial Function

- Many mathematical functions can be defined and solved recursively.
 - The following is a function that computes **n!**.

```
1.  /*
2.   * Compute n! using a recursive definition
3.   * Pre: n >= 0
4.   */
5.  int
6.  factorial(int n)
7.  {
8.      int ans;
9.
10.     if (n == 0)
11.         ans = 1;
12.     else
13.         ans = n * factorial(n - 1);
14.
15.     return (ans);
16. }
```

Iterative factorial Function

- The previous factorial function can also be implemented by a for loop.
 - The **iterative** implementation is usually more efficient than **recursive** implementation.

```
1.  /*
2.   * Computes n!
3.   * Pre: n is greater than or equal to zero
4.   */
5.  int
6.  factorial(int n)
7.  {
8.      int i,          /* local variables */
9.      product = 1;
10.
11.     /* Compute the product n x (n-1) x (n-2) x ... x 2 x 1 */
12.     for (i = n; i > 1; --i) {
13.         product = product * i;
14.     }
15.
16.     /* Return function result */
17.     return (product);
18. }
```

Recursive fibonacci Function

- In our midterm, you are required to write a fibonacci function.
 - It can be easily solved by the recursive function
 - But the recursive function is inefficient because the same fibonacci values may be computed more than once.

```
4.  */
5.  int
6.  fibonacci(int n)
7.  {
8.      int ans;
9.
10.     if (n == 1 || n == 2)
11.         ans = 1;
12.     else
13.         ans = fibonacci(n - 2) + fibonacci(n - 1);
14.
15.     return (ans);
16. }
```

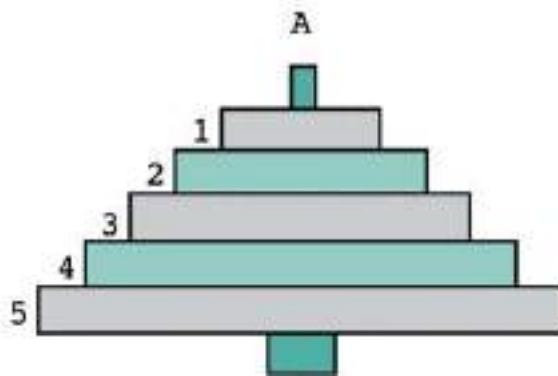
Recursive gcd Function

- Generally speaking, if the algorithm to a problem is defined recursively in itself, we would tend to use the recursive function.
- e.g., the greatest common divisor (GCD) of two integers m and n can be defined recursively.
 - $\text{gcd}(m,n)$ is n if n divides m evenly;
 - $\text{gcd}(m,n)$ is $\text{gcd}(n, \text{remainder of } m \text{ divided by } n)$ otherwise.

```
11. int
12. gcd(int m, int n)
13. {
14.     int ans;
15.
16.     if (m % n == 0)
17.         ans = n;
18.     else
19.         ans = gcd(n, m % n);
20.
21.     return (ans);
22. }
```

A Classical Case: Towers of Hanoi

- The towers of Hanoi problem involves moving a number of disks (in different sizes) from one tower (or called “peg”) to another.
 - The constraint is that the larger disk can never be placed on top of a smaller disk.
 - Only one disk can be moved at each time
 - Assume there are three towers available.



Source

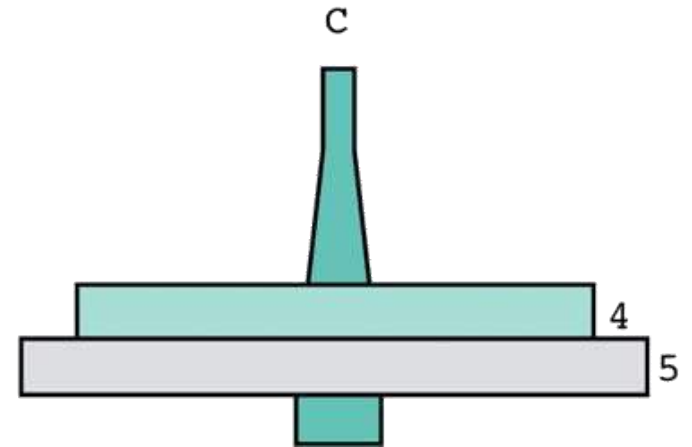
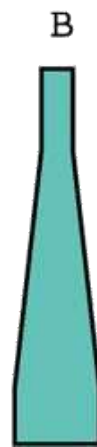
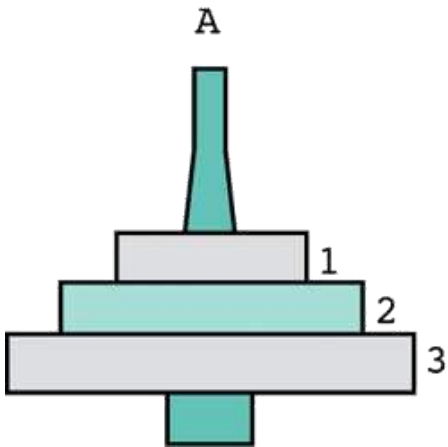
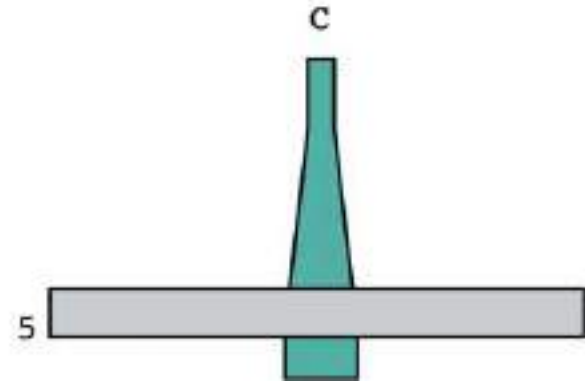
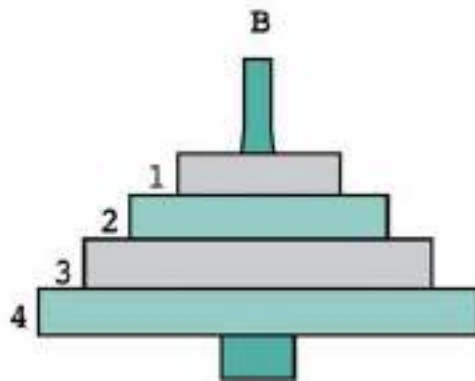


Temp



Destination

A Classical Case: Towers of Hanoi



A Classical Case: Towers of Hanoi

- This problem can be solved easily by recursion.
- Algorithm:

if n is 1 then

 move disk 1 from the source tower to the destination tower

else

1. move $n-1$ disks from the source tower to the temp tower.
2. move disk n from the source tower to the destination tower.
3. move $n-1$ disks from the temp tower to the source tower.

A Classical Case: Towers of Hanoi

```
1.  /*
2.  * Displays instructions for moving n disks from from_peg to to_peg using
3.  * aux_peg as an auxiliary. Disks are numbered 1 to n (smallest to
4.  * largest). Instructions call for moving one disk at a time and never
5.  * require placing a larger disk on top of a smaller one.
6.  */
7.  void
8.  tower(char from_peg,    /* input - characters naming          */
9.        char to_peg,    /*          the problem's
10.       char aux_peg,   /*          three pegs
11.       int n)         /* input - number of disks to move
12.  {
13.     if (n == 1) {
14.         printf("Move disk 1 from peg %c to peg %c\n", from_peg, to_peg);
15.     } else {
16.         tower(from_peg, aux_peg, to_peg, n - 1);
17.         printf("Move disk %d from peg %c to peg %c\n", n, from_peg, to_peg);
18.         tower(aux_peg, to_peg, from_peg, n - 1);
19.     }
20. }
```

The recursive step

The recursive step

A Classical Case: Towers of Hanoi

The execution result of calling `Tower ('A', 'B', 'C', 3);`

Move disk 1 from A to C
Move disk 2 from A to B
Move disk 1 from C to B
Move disk 3 from A to C
Move disk 1 from B to A
Move disk 2 from B to C
Move disk 1 from A to C